THE FUTURE OF COMPUTER SECURITY

Richard Harvey



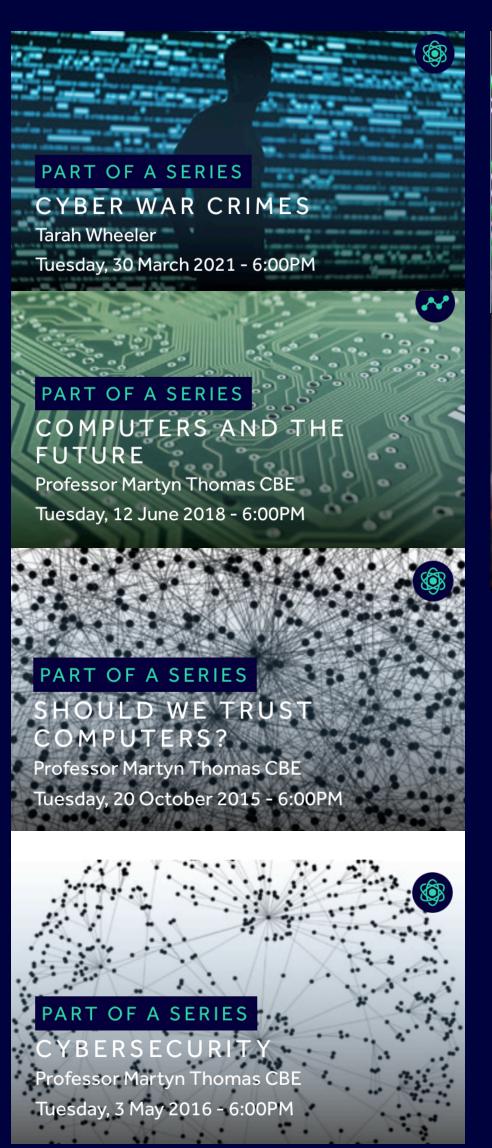
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IT Livery Company Professor of Information Technology, Gresham College



Gresham recap



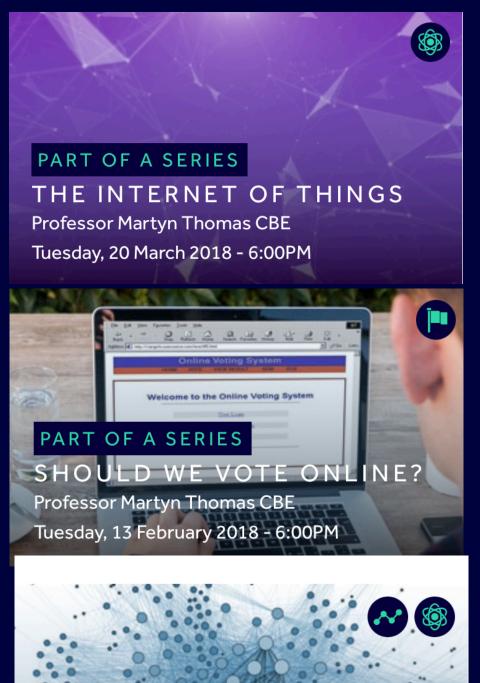


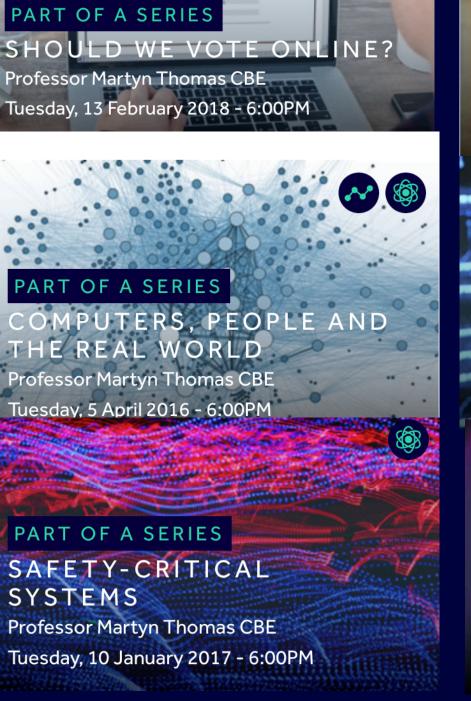














Recap

Developers make mistakes

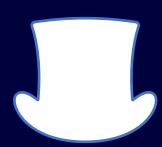


Baddies exploit those mistakes



Goodies, trying to stop other baddies, exploit these mistakes

It's all getting out of hand and innocent people are dying

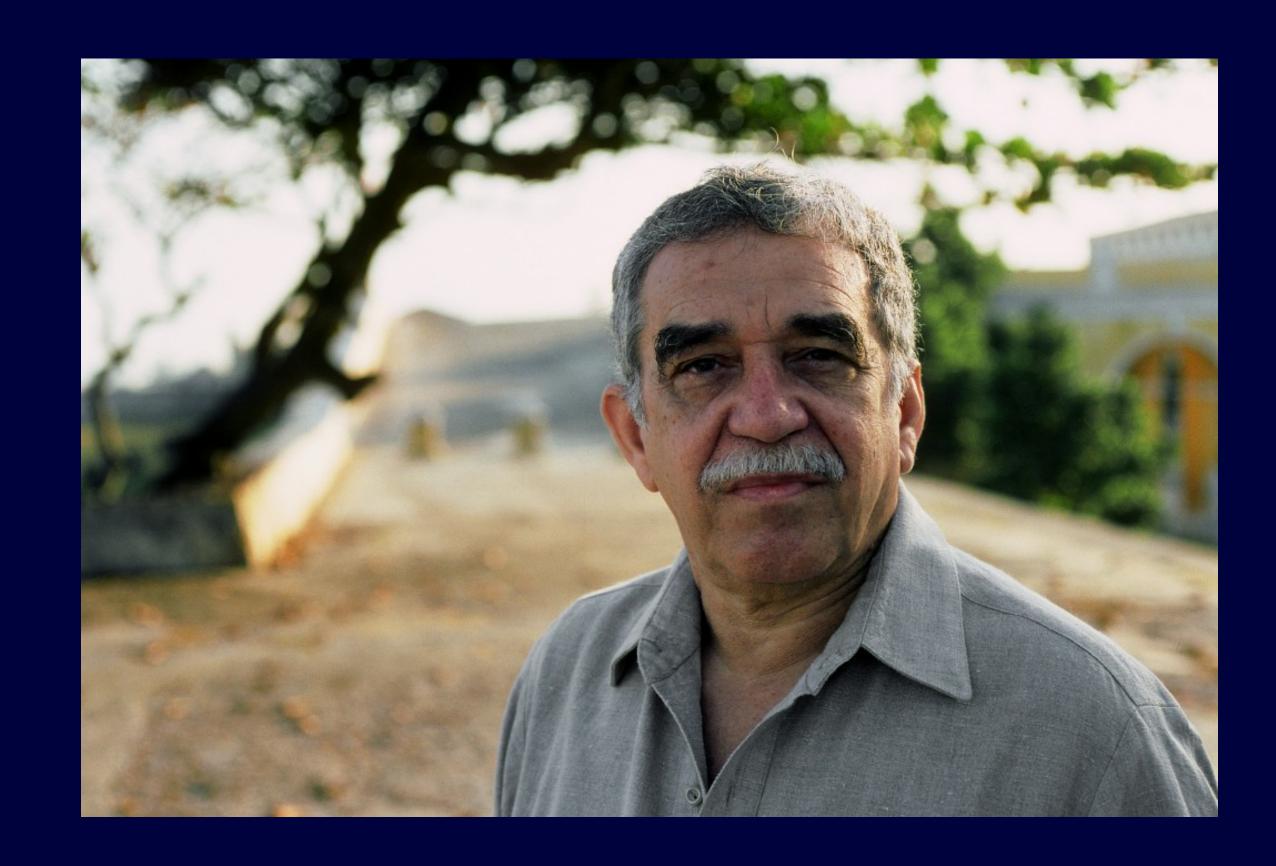


Root cause:

the internet uses unencrypted packets

What this lecture is not about.

"Everyone has three lives: a public life, a private life and a secret life."

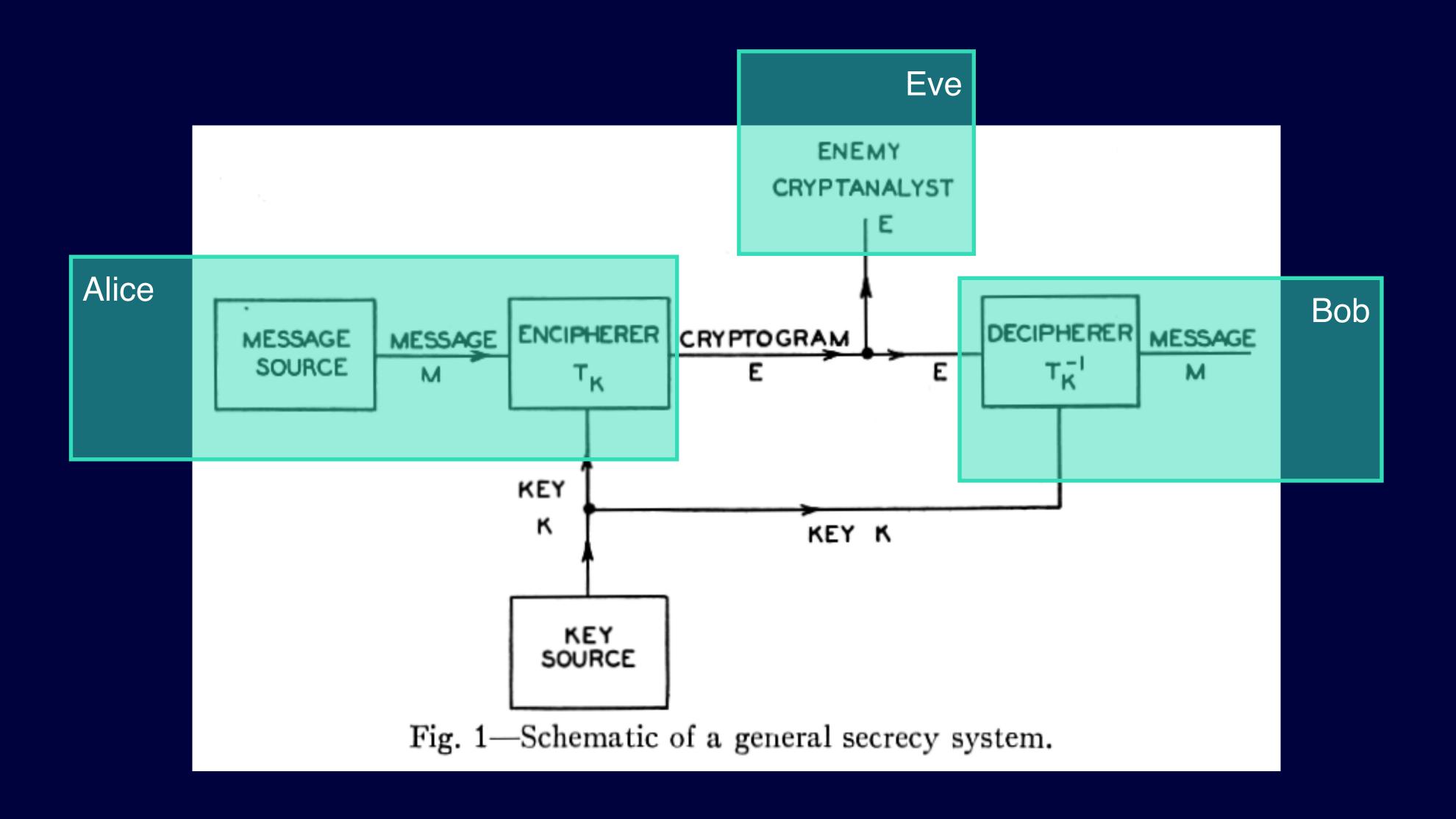


If there was only one form of communication...



Four danger factors

- 1. Internet packets are open
- 2. Internet laws do not mirror in-person laws
- 3. People are disinhibited on the internet
- 4. Birds of a feather can easily flock together



from C. E. Shannon, "Communication theory of secrecy systems," in *The Bell System Technical Journal*, vol. 28, no. 4, pp. 656-715, Oct. 1949

Diffie Hellman Merkle Key Exchange

Basic idea:

Publicise two numbers, G, and, p, and use the following equation to compute public number, K:

 $G^{key} \mod p = K$

What does this mean?

Modulo arithmetic

Let's choose a base, G = 5 $G^1 = 5$; $G^2 = 25$; $G^3 = 125$; $G^4 = 625$...

Now what if we use *modulo* arithmetic, let's choose p = 23.

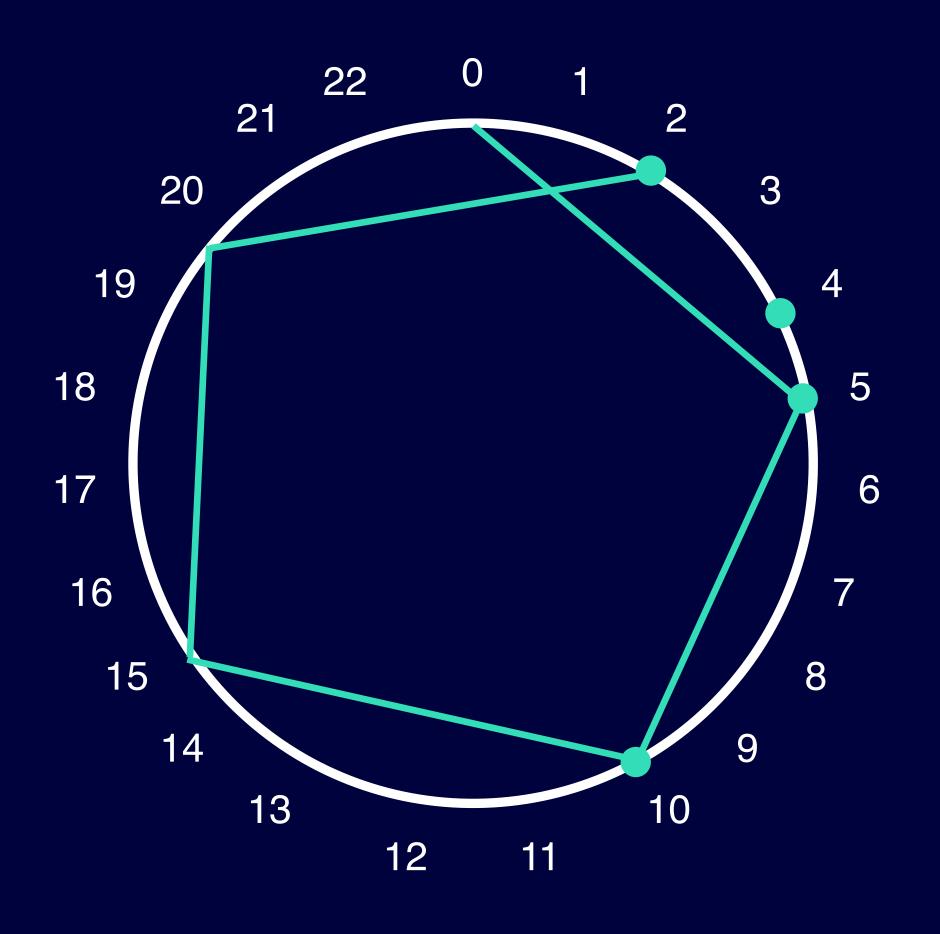
So we divide by p = 23 and compute the remainder

 $G^1 \mod 23 = 5;$

 $G^2 \mod 23 = 2;$

 $G^3 \mod 23 = 10;$

 $G^4 \mod 23 = 4$



Diffie Hellman Merkle Key Exchange

Alice

Secret key a = 6;

Compute public key,

 $A = G^a \mod p = 5^6 \mod 23 = 8$

Compute shared secret key

 $s_a = B^a \mod p = 19^6 \mod 23 = 2$

Public (Eve)

p = 23; G = 5;

A = 8

B = 19

Bob

Secret key b = 15;

Compute public key, $B = 5^{15} \mod 23 = 19$

Compute shared secret key

 $s_b = A^b \mod p = 8^{15} \mod 23 = 2$

 $s_a = B^a \mod p$

 $= (G^b \mod p)^a \mod p$

 $= G^{ab} \mod p = s_b$

How to break Diffie Hellman Merkle

We know the base, G = 5, we know the modulo number p = 23.

Alice has just sent us her encoded key A = 8.

Exhaustive search...

 $5^1 \mod 23 = 5$ No!

 $5^2 \mod 23 = 2$ No!

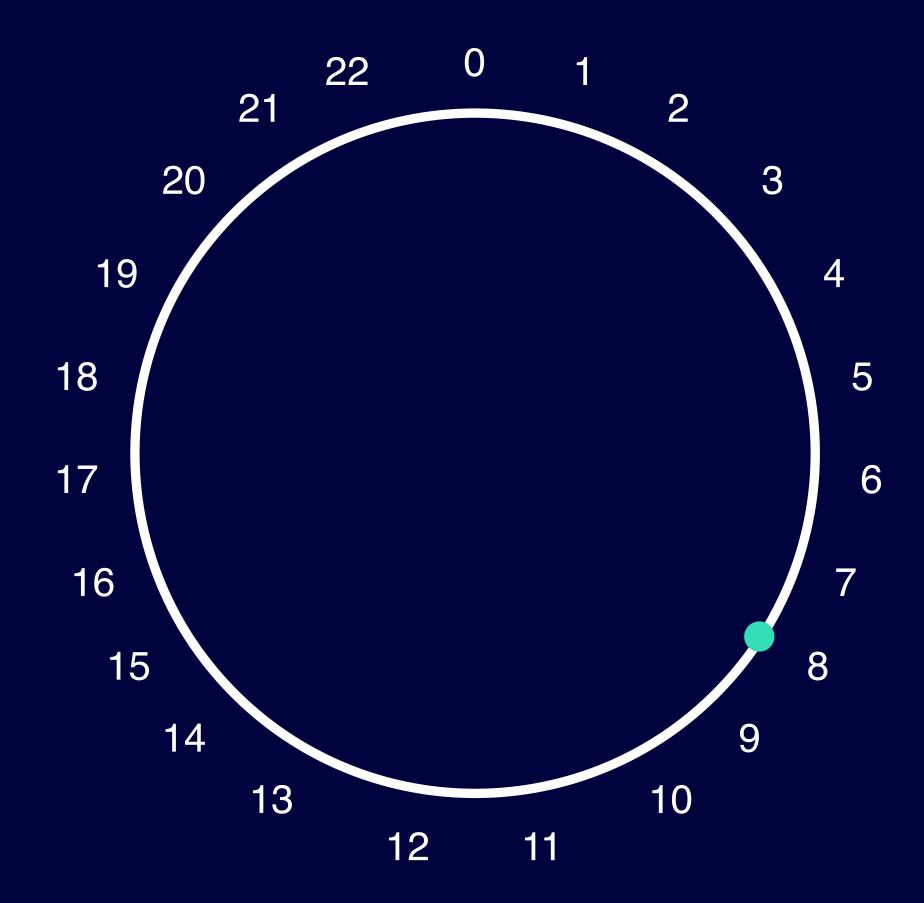
 $5^3 \mod 23 = 10 \text{ No!}$

 $5^4 \mod 23 = 4$ No!

Ad nauseum! We have to check all combinations. If we assume $a \le p$ then we have p checks to make. The qth check involves q multiplications so a full check involves p^2 multiplications.

Note1: if p = 23 then this is not daunting. But a real p might be ...

15525180923007089351309181312584817556313340494345143132023511 94902966239949102107258669453876591642442910007680288864229150 80371891804634263272761303128298374438082089019628850917069131 6593175367469551763119843371637221007210577919 $\sim 10^{464}$ with G = 2



Warning

Do not use G = 2,

1,552,518,092,300,708,935,130,918,131,258,481,755,631,334,049,434,514,313,202,351,194,902,966 ,239,949,102,107,258,669,453,876,591,642,442,910,007,680,288,864,229,150,803,718,918,046,342, 632,727,613,031,282,983,744,380,820,890,196,288,509,170,691,316,593,175,367,469,551,763,119,8 43,371,637,221,007,210,577,919

It has probably already been hacked!

Use a longer p, maybe over 1500 bits long!

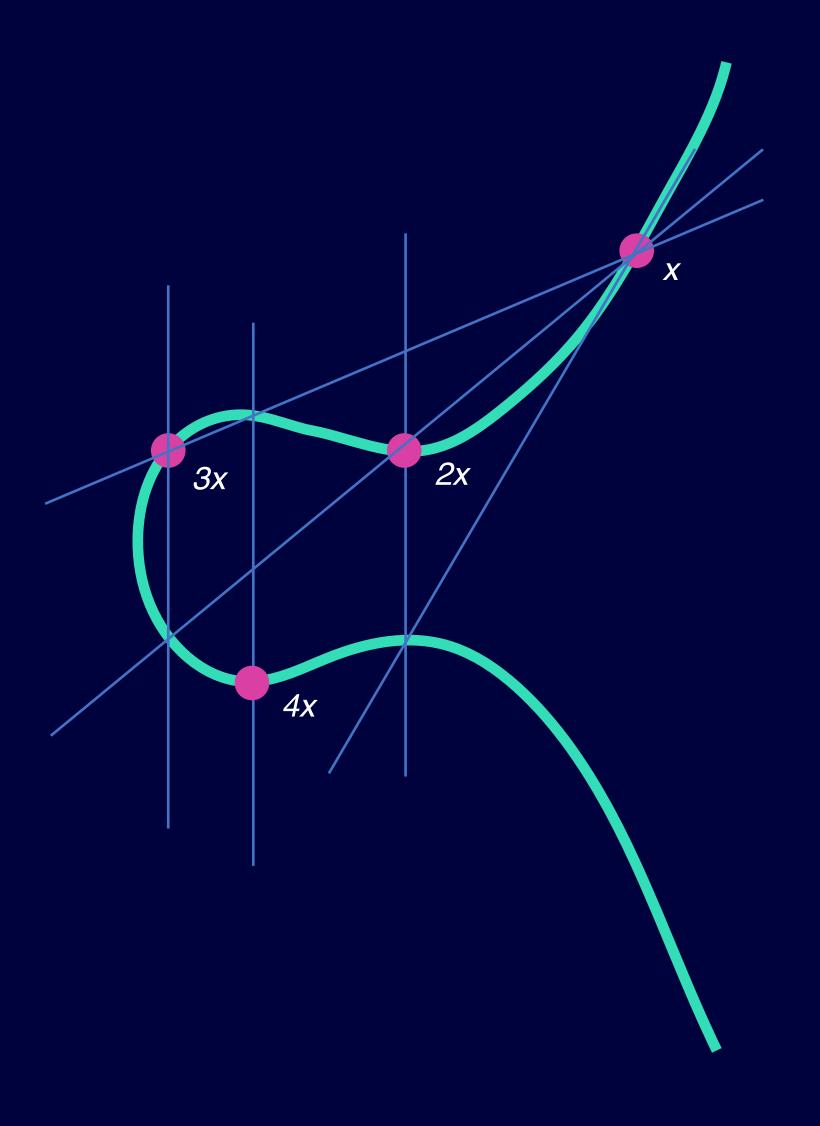
Problem: as the key gets longer and longer encryption takes longer and longer

Solution: instead of using a circle to mix things up (modulo artithmetic) maybe we could use a curve?

Eliptic curves

More complicated than modulo exponentiation so

- shorter keys
- harder to crack
 the "Swiss Army knife of cryptography"



A brief history of hacking

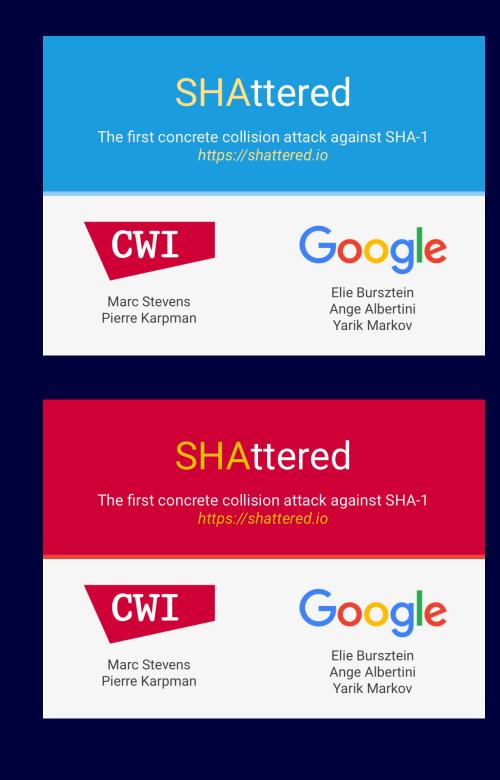
- Leaving aside professional crytography...
 - Early 1980s: phone phreaking
 - 1984: Chaos Computer Club
 - 1988 Morris worm
 - Early 1990s internet used for commerce and everything changes.
 - Three themes: intellectual curiosity; theft and/or extortion and cyber warfare

Hash function

- How can we check that some transmitted data is correct?
- Easy discussed in previous lecture, we use a check number that travels with the data
 - 1. Add a One so that the number of digits in the file is even parity
 - 2. Sum up the Ones and append the count a checksum
 - 3. Generate a string of digits that is a sort of digital finger print for the data.

Cryptographic hash functions

- No inverse
- No collisions
- SHA-1 broken in 2017





38 76 2c f5 59 34 9a e6 a4 c8 ad

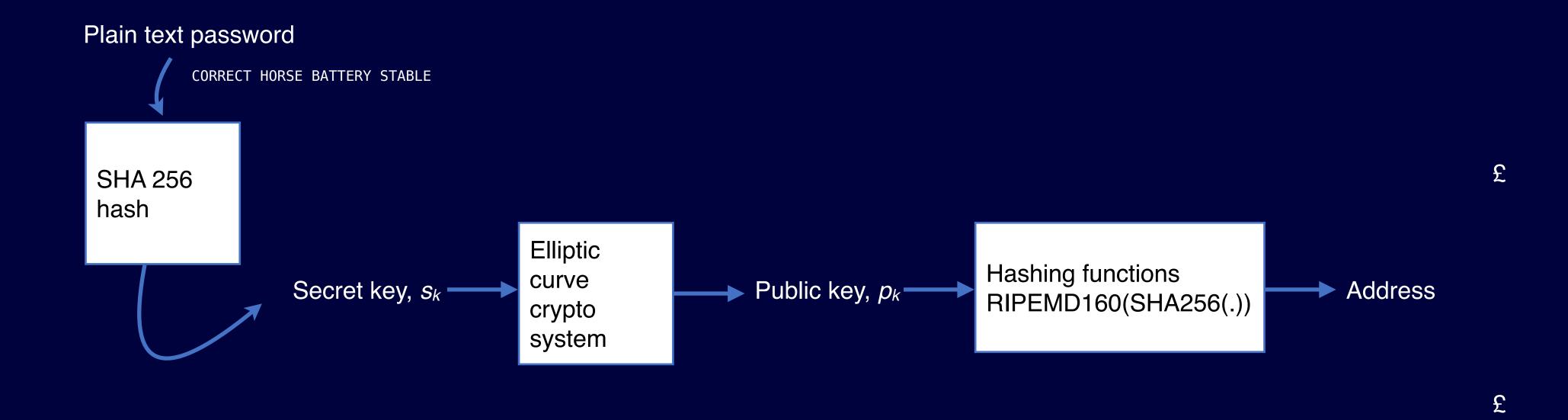
SHA256 Hash function



d6e7c4bdba6e95522b6a3e9263c78d3496306e16662485929e15f1ebc3eb56a1

- 1. Not broken (yet)
- 2. Longer hash
- 3. Still speedy
- 4. Commonly used

Crypto-currency hashes



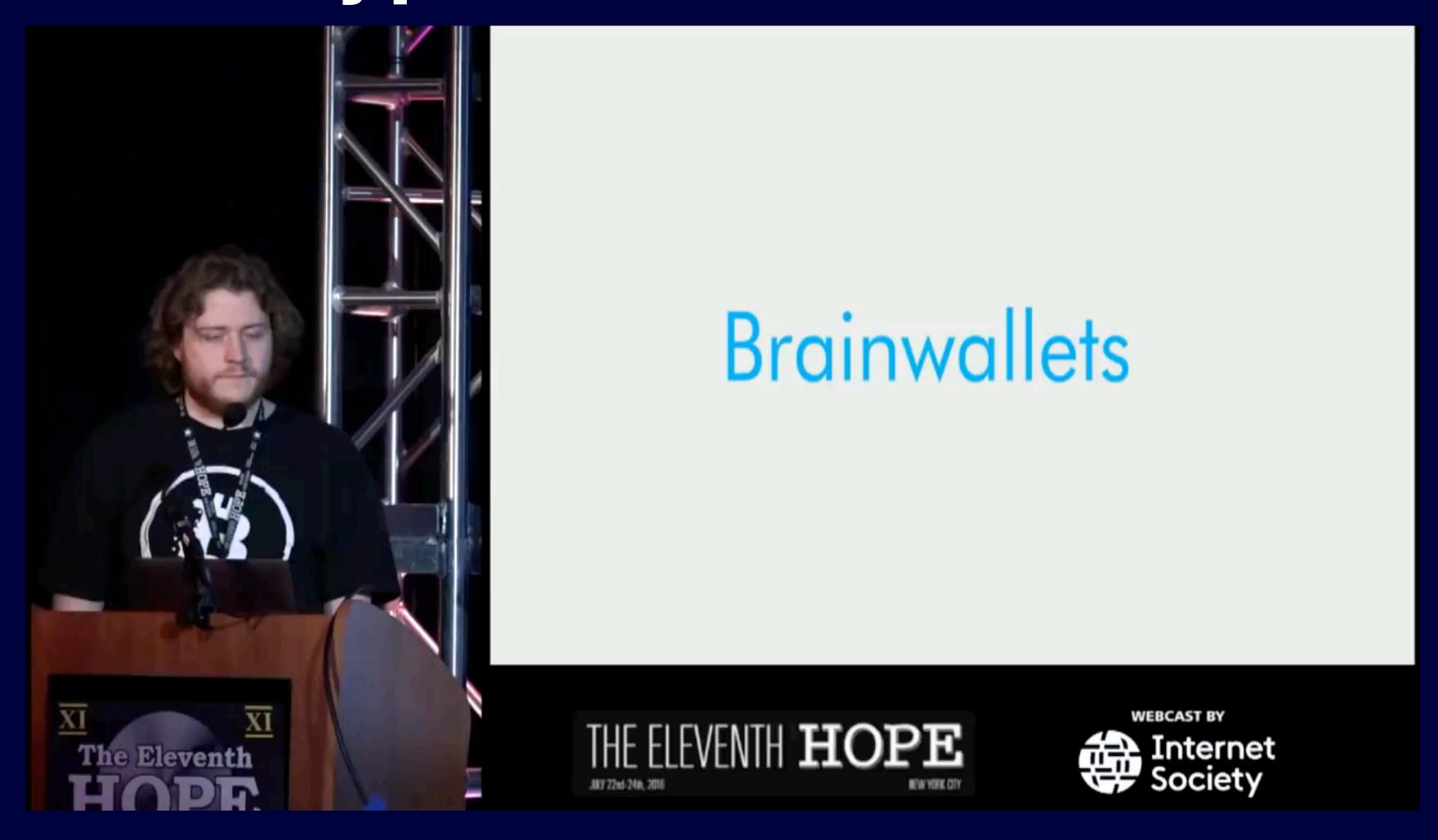
E9873D79C6D87DC0FB6A5778633389_SAMPLE_PRIVATE_KEY_D0_NOT_IMPORT_F4453213303DA61F20BD67FC233AA33262

0450863AD64A87AE8A2FE83C1AF1A8403CB53F53E486D8511DAD8A04887E5B23522CD470243453A299FA9E77237716103ABC11A1DF38855ED6F2EE187E9C582BA6

17VZNX1SN5NtKa8UQFxwQbFeFc3iqRYhem

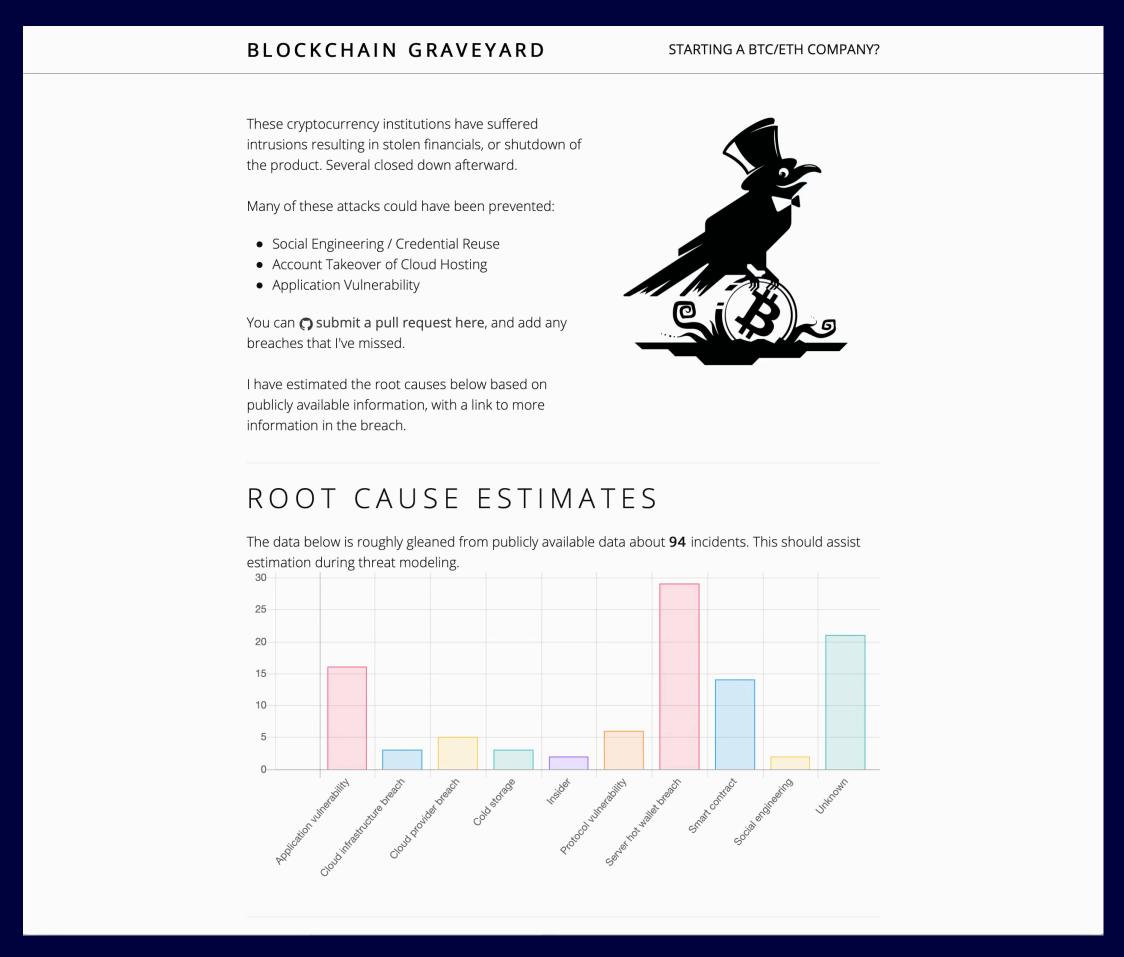


Bad crypto hashes



Public ledger attacks

- Lend themselves to "birthday problem" type attacks
- And each "hit" has a monetary reward



https://magoo.github.io/Blockchain-Graveyard/

State actors

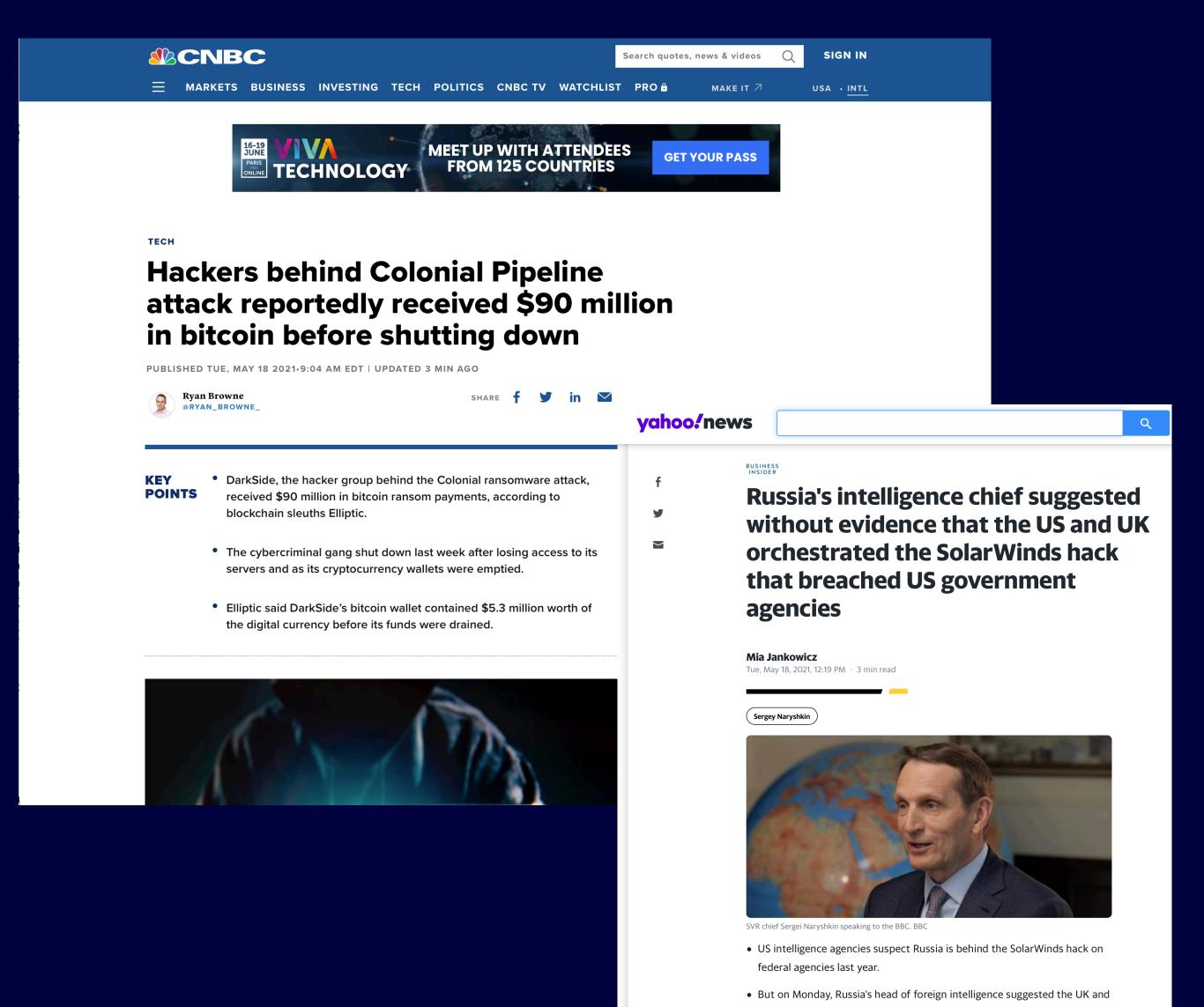
- Snowden leaks describe three types of activity
- 1. Code breaking
- 2. Industrial espionage creation of vulnerabilities
- 3. Hacking collecting and looking

Vulnerabilites Equities Process

- Zero day vulnerabilities vulnerabilities that the manufacturer is unaware of (but is known to someone)
- Should the state hoard knowledge of such vulnerabilities?
- A new twist on the classic ethical dilemma of intelligence gathering...
- Is it possible to gather intelligence that is so valuable that it cannot be revealed?



Last week's news





The role of LEAs

"Sadly, at a time when we need to be taking more action, Facebook are pursuing end-to-end encryption plans that place the good work and progress achieved so far in jeopardy ... This is not acceptable. We cannot allow a situation where law enforcement's ability to tackle abhorrent criminal acts and protect victims is severely hampered."

Pritti Patel, 19th April 2021

"The hack on Solar Winds has shown that state actors have significant capability. We need to be able to understand that threat, protect ourselves from it, and bolster our cyber resilience. While addressing the danger from state and state-sponsored actors, it is rightly a key priority. We also know that criminal groups have the intent and technical means to operate in cyber space."

Pritti Patel speech to CyberUk Conference 11th May 2021



By Richard Townshend - https://members.parliament.uk/member/4066/portrait, CC BY 3.0, https://commons.wikimedia.org/w/index.php?curid=86678234

Quo vadis security?

- Public confusion over privacy rights versus security
- Western governments in denial over strong encryption
- Two recent features:
 - huge financial wins for successful hacks;
 - government sponsored hacking.
- Maybe wait a while before investing in bitcoin?

And...if you fancy getting rich quick...



Next season

Six tech inventions that changed the world October 2021

Thanks and kudos to the Worshipful Company of Information Technologists who sponsor these lectures.

Special thanks to Laurie Mercer from HackerOne for talking me through some of the examples in this talk.